Chapter 17: Distributed Multimedia Systems

- Introduction
- Characteristics of multimedia data
- Quality of service management
- Resource management
- Stream adaptation
- Case study: the Tiger video file server

Introduction

Distributed multimedia applications

- Networked video library, Internet telephony, videoconference
- Generate and consume continuous streams of data in real time

• Characteristics of multimedia applications

- Timely delivery of streams of multimedia data to end-users
 - Audio sample, video frame
- To meet the timing requirements
 - QoS(quality of service)

A typical distributed multimedia system



⇔

The window of scarcity

• A history of computer systems that support distributed data access



Chapter 17: Distributed Multimedia Systems

- Introduction
- Characteristics of multimedia data
- Quality of service management
- Resource management
- Stream adaptation
- Case study: the Tiger video file server

Characteristics of multimedia data

• Continuous

- Refer to the user's view of the data
- Video: a image array is replaced 25 times per second
- Audio: the amplitude value is replaced 8000 times per second
- Time-based
 - The times at which the values are played or recorded affect the validity of the data
 - Hence, the timing should be preserved

Chapter 17: Distributed Multimedia Systems

- Introduction
- Characteristics of multimedia data
- Quality of service management
- Resource management
- Stream adaptation
- Case study: the Tiger video file server

Traditional computer systems

• Multimedia App. compete with other App. for

- Processor cycles, bus cycles, buffer capacity
- Physical transmission links, switches, gateways
- Best-effort policies
 - Multi-task OS: round-robin scheduling or other scheduling
 - shares the processing resource on a best-effort basis among all of the task currently competing
 - Ethernet
 - manages a shared transmission medium in best-efforts manner
- Best-effort policies are not fit to multimedia apps.
- In order to achieve timely delivery,
 - applications need guarantees that the necessary resources will be allocated and scheduled at the required times
 - → QoS management

Typical infrastructure components for multimedia applications



⇔

QoS specifications for components of the application shown in Figure 15.4

Component			Bandwidth	Latency	Loss rate	Resources required
	Camera	Out:	10 frames/sec, raw video 640x480x16 bits) -	Zero	_
A	Codec	In:	10 frames/sec, raw video	Interactive	Low	10 ms CPU each 100 ms;
		Out:	MPEG-1 stream			10 Mbytes RAM
В	Mixer	In:	2×44 kbps audio	Interactive	Very low	1 ms CPU each 100 ms;
		Out:	1 × 44 kbps audio			1 Mbytes RAM
Η	Window	In:	various	Interactive	Low	5 ms CPU each 100 ms;
	system	Out:	50 frame/sec framebuffe	r		5 Mbytes RAM
K	Network connection	In/Out:	MPEG-1 stream, approx 1.5 Mbps	.Interactive	Low	1.5 Mbps, low-loss stream protocol
L	Network connection	In/Out:	Audio 44 kbps	Interactive	Very low	44 kbps, very low-loss stream protocol

The QoS manager's task

- The OoS manager's two main subtasks are:
 - Quality of service negotiation
 - Admission control

The QoS manager's task



Quality of service negotiation

- Resource requirements specification
 - -bandwidth
 - The rate at which a multimedia stream flows
 - -Latency
 - The time required for an individual data element to move through a stream from the source to the destination
 - jitter
 - -Loss rate
 - Data loss due to unmet resource requirements
 - A rate of data loss that can be accepted. E.g., 1%

The usage of resource requirements spec.

- Describe a multimedia stream
 - Describe the characteristics of a multimedia stream in a particular environment
 - E.g. a video conference
 - Bandwidth: 1.5Mbps; delay: 150ms, loss rate: 1%
- Describe the resources
 - Describe the capabilities of resources to transport a stream
 - E.g. a network may provide
 - Bandwidth: 64kbps; delay: 10ms; loss rate: 1/1000

Traffic shaping

• Traffic shaping

- Output buffering to smooth the flow of data elements
- Leaky bucket, Token bucket

• The leaky bucket algorithm

- completely eliminate burst
- R
 - A stream will never flow with a rate higher than R
- B
 - Size of the buffer
 - Bound the time for which an element will remain in the buffer

(a) Leaky bucl

Traffic shaping (2)

• The token bucket algorithm

- Allow larger burst
- Token is generated at a fixed rate of R
- the tokens are collected in a bucket of size B
- Data of size *S* can be sent only if at least *S* tokens are in the bucket
- Ensure: over any interval t, the amount of data sent is not larger than Rt+B,



Flow specification – RFC 1363

• Bandwidth

- The maximum transmission unit and maximum transmission rate
- The burstiness of the stream
 - The token bucket size and rate
- Delay
 - The minimum delay that an application can notice, the maximum jitter it can accept
- Loss rate
 - The total acceptable number of losses over a certain interval
 - The maximum number of consecutive losses

Admission control

- Avoid resource overload
- Protect resource from requests that they cannot fulfill
- Bandwidth reservation
 - Reserve some portion of resource bandwidth exclusively
- Statistical multiplexing
 - Reserve minimum or average bandwidth
 - Handle burst that cause some service drop level occasionally
 - Hypothesis
 - a large number of streams the aggregate bandwidth required remains nearly constant regardless of the bandwidth of individual streams

Chapter 17: Distributed Multimedia Systems

- Introduction
- Characteristics of multimedia data
- Quality of service management
- Resource management
- Stream adaptation
- Case study: the Tiger video file server
- Summary

Resource Management

- To provide a certain QoS level to an application, a system needs to have sufficient resources, it also needs to make the resources available to an application when they are needed (scheduling).
- Resource Scheduling: A process needs to have resources assigned to them according to their priority. Following 2 methods are used:
 - Fair Scheduling
 - Real-time scheduling

Resource scheduling

• Fair scheduling

- when several streams compete for the same resource
- Round-robin
 - Packet-by-packet
 - Bit-by-bit

• Real-time scheduling

- Earliest-deadline-first (EDF)
 - Each media element is assigned a *deadline* by which it must be sent out
 - The scheduler send media elements according to their *deadline*

Chapter 17: Distributed Multimedia Systems

- Introduction
- Characteristics of multimedia data
- Quality of service management
- Resource management
- Stream adaptation
- Case study: the Tiger video file server

Stream adaptation

- Stream adaptation
 - An application adapt to changing QoS levels when a certain QoS cannot be guaranteed
- Drop pieces of information
 - Audio stream
 - Drop can be noticed immediately by the listener
 - Video stream
 - Motion JPEG: easy since frames are independent
 - MPEG: difficult since frames are interdependent
- Increase delay
 - Acceptable for non-interactive applications
- Two methodologies are used:
 - <u>Scaling</u>
 - Filtering

Scaling

- When to perform scaling
 - Adapt a stream to the bandwidth available in the system before it enters a bottleneck resource
- Scaling approach
 - Implementation
 - A monitor process at the target
 - A scaler process at the source
 - Monitor keeps track of the arrival times of messages in a stream. Delayed messages are an indication of bottle neck in the system.
 - Monitor sends a scale-down message to the source that scales up again

Different scaling methods

Temporal scaling

- Decrease the number of video frames transmitted within an interval

Spatial scaling

 Reduce the number of pixels of each image in an video stream, e.g., JPEG and MPEG-2

• Frequency scaling

- Modify the compression algorithm applied to an image

Amplitudinal scaling

- Reduce the color depths for each image pixel

• Color space scaling

 Reduce the number of entries in the color space, e.g., from color to grey-scale presentation

Filtering

- Scaling is not suitable to a stream that involves several receivers
 - Since scaling is conducted at the source, a *scale-down* message will degrade the quality of all streams
- Filtering
 - A stream is partitioned into a set of hierarchical sub-streams
 - The capacity of nodes on a path determines the number of substreams a target receives



Chapter 17: Distributed Multimedia Systems

- Introduction
- Characteristics of multimedia data
- Quality of service management
- Resource management
- Stream adaptation
- Case study: the Tiger video file server

Design goals

- Video-on-demand for a large number of users
 - A large stored digital movie library
 - Users can perform pause, rewind, fast-forward
- Quality of service
 - Constant rate
 - a maximum jitter and low loss rate
- Scalable and distributed
 - Support up to 10000 clients simultaneously
- Low-cost hardware
 - Constructed by commodity PC
- Fault tolerant
 - Tolerant to the failure of any single server or disk

System architecture

One controller

- Connect with each server by low-bandwidth network

• Cubs – the server group

- Each cub is attached by a number of disks (2-4)
- Cubs are connected to clients by ATM



Storage organization

- Stripping
 - A movie is divided into blocks
 - The blocks of a movie are stored on disks attached to different cubs in a sequence of the disk number
 - Deliver a movie: deliver the blocks of the movie from different disks in the sequence number
 - Load-balance when delivering hotspot movies
- Mirroring
 - Each block is divided into several portions (secondaries)
 - The secondaries are stored in the successors
 - If a block is on a disk *i*, then the *secondaries* are stored on disks *i*+1 to *i*+*d*
 - Fault-tolerance for single cub or disk failure

Distributed schedule

- Slot
 - The work to be done to play one block of a movie

• Deliver a stream

- Deliver the blocks of the stream disk by disk
- Can be viewed as a slot moving along disks step by step

• Deliver multiple streams

- Multiple slots moving along disks step by step

• Viewer state

- Address of the client computer
- Identity of the file being played
- Viewer's position in the file
- The viewer's play sequence number
- Bookeeping information

Tiger schedule



Distributed schedule (2)

- Block play time T
 - The time that will be required for a viewer to display a block on the client computer
 - Typically about 1 second for all streams
 - The next block of a stream must begin to be delivered *T* time after the current block begin to be delivered
- *Block service time t* (a slot)
 - Read the next block into buffer
 - Deliver it to the client
 - Update viewer state in the schedule and pass the updated slot to the next cub
 - T / t typically result in a value > 4
- The maximum streams the Tiger system can support simultaneously
 - T/t * the number of disks